

# Quantum expanders – Random constructions & Applications

Based on the works:

- arXiv:1906.11682 (with David Pérez-García)
- arXiv:2302.07772 (with Pierre Youssef)
  - arXiv:2409.17971

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- 1 Classical and quantum expanders: definitions and motivations
- 2 Random constructions of expanders
- 3 Applications and perspectives

$G$  a (directed or undirected)  $d$ -biregular graph on  $n$  vertices.

↳  $d$  incoming and  $d$  outgoing edges at each vertex

$A$  its (normalized) adjacency matrix, i.e. the  $n \times n$  matrix s.t.  $A_{kl} = e(l \rightarrow k)/d$  for all  $1 \leq k, l \leq n$ .  
number of edges from vertex  $l$  to vertex  $k$  ↵

$\lambda_1(A), \dots, \lambda_n(A)$  eigenvalues of  $A$ , ordered s.t.  $|\lambda_1(A)| \geq \dots \geq |\lambda_n(A)|$ .

$G$  biregular  $\Rightarrow A$  bistochastic  $\Rightarrow \lambda_1(A) = 1$ , with associated eigenvector the uniform probability  $u$ .  
The *spectral expansion parameter* of  $G$  is  $\lambda(G) := |\lambda_2(A)|$ .  $(1/n, \dots, 1/n) = \leftarrow$

**Observation:**  $\lambda(G) = |\lambda_1(A - J)|$ , where  $J$  is the adjacency matrix of the *complete graph* on  $n$  vertices, i.e. the matrix whose entries are all equal to  $1/n$ .

$\rightarrow \lambda(G)$  is a distance measure between  $G$  and the complete graph.

# Classical expanders

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## Definition [Classical expander (informal)]

A  $d$ -biregular graph  $G$  on  $n$  vertices is an *expander* if it is sparse (i.e.  $d \ll n$ ) and spectrally expanding (i.e.  $\lambda(G) \ll 1$ ).

$\rightarrow G$  is both 'economical' and 'resembling' the complete graph.

For instance, a random walk supported on  $G$  converges fast to equilibrium:

For any probability  $p$  on  $\{1, \dots, n\}$ ,  $\forall q \in \mathbf{N}$ ,  $\|A^q p - u\|_1 \leq \sqrt{n} \|A^q p - u\|_2 \leq \sqrt{n} \lambda(G)^q$ .

exponential convergence, at rate  $|\log \lambda(G)|$  ↵

# Quantum analogue of the transition matrix associated to a biregular graph

Classical - Quantum correspondence:

- $p \in \mathbf{R}^n$  probability vector  $\longleftrightarrow \rho \in \mathcal{M}_n(\mathbf{C})$  density operator (aka *quantum state*).  
↳ self-adjoint positive semidefinite trace 1 operator
- $A : \mathbf{R}^n \rightarrow \mathbf{R}^n$  transition matrix  $\longleftrightarrow \Phi : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C})$  *quantum channel*.  
↳ completely positive (CP) trace-preserving (TP) linear map
- $G$  biregular:  $A$  leaves  $u$  invariant  $\longleftrightarrow \Phi$  unital:  $\Phi$  leaves  $I/n$  invariant.  
↳ maximally mixed density operator

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 $\hookrightarrow$  maximally mixed density operator

**Question:** What is the analogue of the degree in the quantum setting?

**Answer:** The Kraus rank.

Given a CP map  $\Phi$  on  $\mathcal{M}_n(\mathbf{C})$ , a *Kraus representation* of  $\Phi$  is of the form:

$$\Phi : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \sum_{i=1}^d K_i X K_i^* \in \mathcal{M}_n(\mathbf{C}), \text{ where } K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C}). \quad (\star)$$

$\hookrightarrow$  Kraus operators of  $\Phi$

The minimal  $d$  s.t.  $\Phi$  can be written as  $(\star)$  is the *Kraus rank* of  $\Phi$  (it is always at most  $n^2$ ).

[ Note:  $\Phi$  is TP iff  $\sum_{i=1}^d K_i^* K_i = I$ .  $\Phi$  is unital iff  $\sum_{i=1}^d K_i K_i^* = I$ . ]

Indeed, the degree and the Kraus rank both quantify the 1-iteration spreading:

- $G$  a degree  $d$  graph: If  $|\text{supp}(p)| = 1$ , then  $|\text{supp}(Ap)| \leq d$ .
- $\Phi$  a Kraus rank  $d$  quantum channel: If  $\text{rank}(\rho) = 1$ , then  $\text{rank}(\Phi(\rho)) \leq d$ .

$\Phi$  a Kraus rank  $d$  unital quantum channel on  $\mathcal{M}_n(\mathbf{C})$ .

$\lambda_1(\Phi), \dots, \lambda_{n^2}(\Phi)$  eigenvalues of  $\Phi$ , ordered s.t.  $|\lambda_1(\Phi)| \geq \dots \geq |\lambda_{n^2}(\Phi)|$ .

$\Phi$  TP and unital  $\Rightarrow \lambda_1(\Phi) = 1$ , with associated eigenstate the maximally mixed state  $I/n$ .

The *spectral expansion parameter* of  $\Phi$  is  $\lambda(\Phi) := |\lambda_2(\Phi)|$ .

**Observation:**  $\lambda(\Phi) = |\lambda_1(\Phi - \Pi)|$ , where  $\Pi$  is the *maximally mixing channel* on  $\mathcal{M}_n(\mathbf{C})$ , i.e.

$\Pi : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \text{Tr}(X) I/n \in \mathcal{M}_n(\mathbf{C})$ .

$\longrightarrow \lambda(\Phi)$  is a distance measure between  $\Phi$  and the maximally mixing channel.

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For instance, the dynamics associated to  $\Phi$  converges fast to equilibrium:

For any state  $\rho$  on  $\mathbf{C}^n$ ,  $\forall q \in \mathbf{N}$ ,  $\|\Phi^q(\rho) - I/n\|_1 \leq \sqrt{n} \|\Phi^q(\rho) - I/n\|_2 \leq \sqrt{n} \lambda(\Phi)^q$ .

exponential convergence, at rate  $|\log \lambda(\Phi)| \nwarrow$



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# Constructions of optimal classical expanders

**Fact:** For any undirected  $d$ -regular graph  $G$ ,  $\lambda(G) \geq 2\sqrt{d-1}/d - o_n(1)$ .

→  $G$  is an optimal undirected expander (aka *Ramanujan graph*) if  $\lambda(G) \leq 2\sqrt{d-1}/d$ .

**Question:** Do Ramanujan graphs exist?

- ① Explicit constructions of exactly Ramanujan graphs only for  $d = p^m + 1$ ,  $p$  prime.
- ② Random constructions of almost Ramanujan graphs for all  $d$ .
- ③ Existence of exactly Ramanujan graphs for all  $d$ .

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- 3 Existence of exactly Ramanujan graphs for all  $d$ .

In fact, for large  $n$ , almost all undirected regular graphs are almost Ramanujan:

**Theorem [Uniform random undirected regular graph (Friedman, Bordenave)]**

Fix  $d \in \mathbf{N}$ . Let  $G$  be uniformly distributed on the set of undirected  $d$ -regular graphs on  $n$  vertices.

Then, for all  $\varepsilon > 0$ ,  $\mathbf{P}\left(\lambda(G) \leq \frac{2\sqrt{d-1}}{d} + \varepsilon\right) = 1 - o_n(1)$ .

**Remarks:**

→ simpler model of random undirected regular graph

- First proven for the ‘doubled edges permutation model’: for  $d \in \mathbf{N}$  even, pick  $\sigma_1, \dots, \sigma_{d/2} \in S_n$  independent uniformly distributed and let  $G$  have edges  $\{(k, \sigma_i(k)), (k, \sigma_i^{-1}(k))\}_{1 \leq k \leq n, 1 \leq i \leq d/2}$ .
- Result remains true for  $G$  a random directed regular graph and for  $d$  growing with  $n$ , up to a constant multiplicative factor:  $\mathbf{P}(\lambda(G) \leq C/\sqrt{d} + \varepsilon) = 1 - o_n(1)$ .

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**Idea:** Pick  $K_1, \dots, K_d \in \mathcal{M}_n(\mathbf{C})$  at random, under the constraints 
$$\begin{cases} \sum_{i=1}^d K_i^* K_i = I \\ \sum_{i=1}^d K_i K_i^* = I \end{cases}.$$

$\Phi : X \mapsto \sum_{i=1}^d K_i X K_i^*$  is a random Kraus rank (at most)  $d$  unital quantum channel on  $\mathcal{M}_n(\mathbf{C})$ .

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## Theorem [Paired Haar unitaries as Kraus operators (Hastings)]

Fix  $d \in \mathbf{N}$  even. Pick  $U_1, \dots, U_{d/2} \in \mathcal{M}_n(\mathbf{C})$  independent Haar unitaries. Let  $K_i = \frac{U_i}{\sqrt{d}}$ ,  $1 \leq i \leq \frac{d}{2}$ .

The self-adjoint unital quantum channel  $\Phi$  associated to the  $K_i$ 's,  $K_i^*$ 's satisfies:

$$\forall \varepsilon > 0, \mathbf{P} \left( \lambda(\Phi) \leq \frac{2\sqrt{d-1}}{d} + \varepsilon \right) = 1 - o_n(1).$$

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**Question:** Can this result be extended to non-self-adjoint random models? to a wider variety of them? to a regime where  $d$  is not fixed but grows with  $n$ ?

# Constructions of optimal non-self-adjoint quantum expanders

**Fact:** For any Kraus rank  $d$  unital quantum channel  $\Phi$ ,  $\lambda(\Phi) \geq 1/\sqrt{d} - o_n(1)$ .

**Theorem [Haar unitaries as Kraus operators (Pisier, Timhadjelt)]**

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The unital quantum channel  $\Phi$  associated to the  $K_i$ 's satisfies:

$$\forall \varepsilon > 0, \mathbf{P} \left( \lambda(\Phi) \leq \frac{1}{\sqrt{d}}(1 + \varepsilon) \right) \geq 1 - e^{-c\varepsilon n^{1/12}}.$$



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**Question:** Does this remain true for 'less random' unitary Kraus operators?

A probability measure  $\mu$  on  $U(n)$  is a  $k$ -design if  $\mathbf{E}_{U \sim \mu} [U^{\otimes k}(\cdot) U^{*\otimes k}] = \mathbf{E}_{U \sim \mu_H} [U^{\otimes k}(\cdot) U^{*\otimes k}]$ .  
Haar measure on  $U(n)$   $\nwarrow$

## Theorem [2-design unitaries as Kraus operators (Lancien)]

Pick  $U_1, \dots, U_d \in \mathcal{M}_n(\mathbf{C})$  independent 2-design unitaries. Let  $K_i = \frac{U_i}{\sqrt{d}}$ ,  $1 \leq i \leq d$ .

If  $d \geq (\log n)^{8+\delta}$ , the unital quantum channel  $\Phi$  associated to the  $K_i$ 's satisfies:

$$\mathbf{P} \left( \lambda(\Phi) \leq \frac{2}{\sqrt{d}} \left( 1 + \frac{C}{(\log n)^{\delta/6}} \right) \right) \geq 1 - \frac{1}{n}.$$

**Interest:** Almost optimal quantum expander from random Kraus operators which are sampled according to a simple measure on the unitary group (uniform measure on explicit finite subset).

## More random examples of optimal quantum expanders

**Need:** Extend the notion of expander to non-unital quantum channels (or else, too constraining).

### Theorem [Gaussians as Kraus operators (Lancien/Pérez-García)]

Pick  $G_1, \dots, G_d \in \mathcal{M}_n(\mathbf{C})$  independent Gaussians. Let  $\tilde{K}_i = \frac{G_i}{\sqrt{d}}$ ,  $K_i = \tilde{K}_i \Sigma^{-1/2}$ ,  $1 \leq i \leq d$ .  
i.i.d. Gaussian entries (mean 0, variance  $1/n$ )  $\swarrow$   $\searrow = \sum_{j=1}^d \tilde{K}_j^* \tilde{K}_j$

The quantum channel  $\Phi$  associated to the  $K_i$ 's satisfies:  $\mathbf{P} \left( \lambda(\Phi) \leq \frac{c}{\sqrt{d}} \right) \geq 1 - e^{-cn}$ .

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**Question:** Does this remain true for 'less random' Kraus operators with independent entries?

### Theorem [Sparse random matrices as Kraus operators (Lancien/Youssef)]

$\rightarrow$  e.g. adjacency matrix of  $d$ -biregular graph  $G$  on  $n$  vertices s.t.  $\lambda(G) \leq C/\sqrt{d}$   
Fix  $A \in \mathcal{M}_n(\mathbf{R})$  a bistochastic matrix s.t.  $|\lambda_2(A)| \leq \frac{C}{\sqrt{d}}$ . Let  $W \in \mathcal{M}_n(\mathbf{C})$  be a random matrix with independent centered entries s.t.  $\forall 1 \leq k, l \leq n$ ,  $\mathbf{E}|W_{kl}|^2 = A_{kl}$ ,  $(\mathbf{E}|W_{kl}|^p)^{1/p} \leq C' p^\beta A_{kl}$ ,  $p \in \mathbf{N}$ .  
[  $\beta = 0$ : bounded,  $\beta = 1/2$ : sub-Gaussian,  $\beta = 1$ : sub-exponential ]  $\swarrow$   
Pick  $W_1, \dots, W_d \in \mathcal{M}_n(\mathbf{C})$  independent copies of  $W$ . Let  $\tilde{K}_i = \frac{W_i}{\sqrt{d}}$ ,  $K_i = \tilde{K}_i \Sigma^{-1/2}$ ,  $1 \leq i \leq d$ .  
If  $d \geq (\log n)^8$ , the quantum channel  $\Phi$  associated to the  $K_i$ 's satisfies:  $\mathbf{P} \left( \lambda(\Phi) \leq \frac{C''}{\sqrt{d}} \right) \geq 1 - \frac{1}{n}$ .

**Interest:** Almost optimal quantum expanders from random Kraus operators which are sparse and whose entries have any distribution following the moments' growth assumption.

## Proof idea to show that $\mathbf{E}\lambda(\Phi) \leq C/\sqrt{d}$

**Goal:** In all cases, we want to upper bound  $\mathbf{E}|\lambda_2(\Phi)| = \mathbf{E}|\lambda_1(\Phi - \Pi_{\rho_*})|$ , where  $\rho_*$  is the fixed state of  $\Phi$  and  $\Pi_{\rho_*} : X \in \mathcal{M}_n(\mathbf{C}) \mapsto \text{Tr}(X)\rho_* \in \mathcal{M}_n(\mathbf{C})$ .

First step: Upper bound  $\mathbf{E}|\lambda_1(\Phi - \mathbf{E}(\Phi))|$  (and then show that  $\mathbf{E}(\Phi)$  is close to  $\Pi_{\rho_*}$ ).

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- Observation 1:  $|\lambda_1(\Psi)| \leq s_1(\Psi) = \|\Psi\|_\infty$ .

- Observation 2:  $\|\Psi\|_\infty = \|M_\Psi\|_\infty$ , where for  $\Psi : X \mapsto \sum_{i=1}^d K_i X L_i^*$ ,  $M_\Psi = \sum_{i=1}^d K_i \otimes \bar{L}_i$ .

[ Identification  $\Psi : \mathcal{M}_n(\mathbf{C}) \rightarrow \mathcal{M}_n(\mathbf{C}) \equiv M_\Psi : \mathbf{C}^n \otimes \mathbf{C}^n \rightarrow \mathbf{C}^n \otimes \mathbf{C}^n$  preserves the operator norm. ]

→ We want to upper bound  $\mathbf{E}\|\underbrace{M_\Phi - \mathbf{E}(M_\Phi)}_{=:X}\|_\infty$ , where  $M_\Phi = \sum_{i=1}^d K_i \otimes \bar{K}_i$  with the  $K_i$ 's random.

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→ Haar unitaries, Gaussians

- For concrete models, this can be done by a moments' method:

By Jensen's inequality, we have:  $\forall p \in \mathbf{N}$ ,  $\mathbf{E}\|X\|_\infty \leq \mathbf{E}\|X\|_p \leq (\mathbf{E}\text{Tr}|X|^p)^{1/p}$ .

The term on the r.h.s. can be estimated and provides a good upper bound for  $p \simeq n^\gamma$ .

↳ by Weingarten or Wick calculus

→ 2-design unitaries, arbitrary sparse random matrices

- For more general cases, we use recent operator norm estimates for random matrices with dependencies and non-homogeneity (Brailovskaya/van Handel):

Setting  $X = \sum_{i=1}^d Z_i$ , with  $Z_i := K_i \otimes \bar{K}_i - \mathbf{E}(K_i \otimes \bar{K}_i)$ ,  $1 \leq i \leq d$ , we have:

$$\mathbf{E}\|X\|_\infty \lesssim \|\mathbf{E}(XX^*)\|_\infty^{1/2} + \|\mathbf{E}(X^*X)\|_\infty^{1/2} + (\log n)^{3/2} \|\mathbf{Cov}(X)\|_\infty^{1/2} + (\log n)^6 \left( \mathbf{E} \max_{1 \leq i \leq d} \|Z_i\|_\infty^2 \right)^{1/2}.$$

- 1 Classical and quantum expanders: definitions and motivations
- 2 Random constructions of expanders
- 3 Applications and perspectives**

# Implications for typical decay of correlations in many-body quantum systems

*Matrix product states (MPS)* form a subset of *many-body quantum states*.

They are particularly useful because:

- They admit an *efficient description*: number of parameters that scales linearly rather than exponentially with the number of subsystems.
- They are *good approximations of several 'physically relevant' states*, such as ground states of gapped local Hamiltonians on 1D systems (Hastings, Landau/Vazirani/Vidick).
  - ↳ composed of terms which act non-trivially only on nearby sites
  - ↳ spectral gap lower bounded by a constant independent of the number of subsystems



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└─ composed of terms which act non-trivially only on nearby sites  
└─ spectral gap lower bounded by a constant independent of the number of subsystems

between observables measured on distinct sites ←  
with the distance separating the sites ←

**Fact:** Random (translation-invariant) MPS typically have correlations that decay exponentially fast, with a *small correlation length* (Lancien/Pérez-García).

**Proof strategy:** Observe that the correlation length is given by  $1/|\log \lambda(\Phi)|$  for  $\Phi$  a random quantum channel associated to the random MPS (its so-called *transfer operator*).

## Some perspectives

- What about *explicit constructions* of optimal quantum expanders?  
Important for applications (cryptography, error correction, condensed matter physics, etc).  
Seminal constructions required a large amount of randomness.  
First step towards *derandomization*: Kraus operators sampled from an explicit finite subset of unitaries or as sparse matrices with independent  $\pm 1$  entries.

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- What about identifying the *full spectral distribution* of random quantum channels?

$\rho_*$  fixed state of  $\Phi$ ,  $\Pi_{\rho_*} : X \mapsto \text{Tr}(X)\rho_*$  ↗

For a random Kraus rank  $d$  quantum channel  $\Phi$  on  $\mathcal{M}_n(\mathbf{C})$ , the eigenvalues of  $\Phi - \Pi_{\rho_*}$  are typically inside a segment of half-length  $2\sqrt{d-1}/d$  or a disc of radius  $1/\sqrt{d}$  for large  $n$ .

↳ self-adjoint case

↳ non-self-adjoint case

But what is their exact distribution inside this segment or disc?

Answer: As  $n, d \rightarrow \infty$ , the spectral distribution of  $\sqrt{d}(\Phi - \Pi_{\rho_*})$  is proven to be semi-circular in the self-adjoint case (Lancien/Oliveira Santos/Youssef) and conjectured to be circular in the non-self-adjoint case (Bruzda/Cappellini/Sommers/Życzkowski, Aubrun/Nechita).

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- Do the results about the typical spectral gap of random quantum channels remain true when we impose *extra symmetries* on the model?
- What about looking at other, related, notions of expansions, such as *geometric* ones (Bannink/Briët/Labib/Maassen) or *linear-algebraic* ones (Li/Qiao/Wigderson/Wigderson/Zhang)?

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